Alternatives for Scripting in Conjunction with an Algebraic Modeling Language for Optimization

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25th European Conference on Operational Research
Vilnius, Lithuania — 8-11 July 2012
Session WD-33, Optimization Modeling II
Alternatives for Scripting in Conjunction with an Algebraic Modeling Language for Optimization

Optimization modeling languages are fundamentally declarative, yet successful languages also offer ways to write scripts or programs. What can scripting in a modeling language offer in comparison to modeling in a general-purpose scripting language? Some answers will be suggested through diverse examples in which the AMPL modeling language is applied to parametric analysis, solution generation, heuristic optimization, pattern enumeration, and decomposition. Concluding comments will touch on the complexity of scripts seen in practical applications, and on prospects for further improvements.
Alternatives for Programming in conjunction with an Algebraic Modeling Language for Optimization

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INFORMS National Meeting
New Orleans, October 30, 1995
Topics: Introduction to AMPL

The optimization modeling cycle

Optimization modeling languages

Example: multicommodity transportation
- Mathematical formulation
- AMPL formulation
- AMPL solution
Topics: Scripting in AMPL

1: *Parametric analysis*

2: *Solution generation*
   a: *via cuts*
   b: *via solver*

3: *Heuristic optimization*

4: *Pattern generation*

5: *Decomposition*

*Scripts in practice . . .*

*Prospective improvements . . .*
The Optimization Modeling Cycle

Steps

- Communicate with problem owner
- Build model
- Prepare data
- Generate optimization problem
- Submit problem to solver
  * CPLEX, Gurobi, KNITRO, CONOPT, MINOS, . . .
- Report & analyze results
- Repeat!

Goals

- Do this quickly and reliably
- Get results before client loses interest
- Deploy for application
What Makes This Hard?

“We do not feel that the linear programming user’s most pressing need over the next few years is for a new optimizer that runs twice as fast on a machine that costs half as much (although this will probably happen). Cost of optimization is just not the dominant barrier to LP model implementation.

“The process required to manage the data, formulate and build the model, report on and analyze the results costs far more, and is much more of a barrier to effective use of LP, than the cost/performance of the optimizer.”

Krabek, Sjoquist, Sommer,
“The APEX Systems: Past and Future.”
Optimization Modeling Languages

Two forms of an optimization problem
- Modeler’s form
  * Mathematical description, easy for people to work with
- Algorithm’s form
  * Explicit data structure, easy for solvers to compute with

Idea of a modeling language
- A computer-readable modeler’s form
  * You write optimization problems in a modeling language
  * Computers translate to algorithm’s form for solution

Advantages of a modeling language
- Faster modeling cycles
- More reliable modeling and maintenance
Algebraic Modeling Languages

Formulation concept

- Define data in terms of sets & parameters
  - Analogous to database keys & records
- Define decision variables
- Minimize or maximize a function of decision variables
- Subject to equations or inequalities that constrain the values of the variables

Advantages

- Familiar
- Powerful
- Implemented
The AMPL Modeling Language

Features
- Algebraic modeling language
- Variety of data sources
- Connections to all solver features
- Interactive and scripted control

Advantages
- Powerful, general expressions
- Natural, easy-to-learn design
- Efficient processing scales well with problem size
AMPL with Gurobi

Features

- Detection of all supported problem types
- Access to all algorithm & display options

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>aggfill</td>
<td>amount of fill allowed during aggregation during Gurobi's presolve (default 10)</td>
</tr>
<tr>
<td>aggregate</td>
<td>whether to use aggregation during Gurobi presolve: 0 = no (sometimes reduces numerical errors), 1 = yes (default)</td>
</tr>
<tr>
<td>ams_eps</td>
<td>relative tolerance for reporting alternate MIP solutions (default = no limit)</td>
</tr>
<tr>
<td>ams_epsabs</td>
<td>absolute tolerance for reporting alternate MIP solutions (default = no limit)</td>
</tr>
<tr>
<td>ams_limit</td>
<td>limit on number of alternate MIP solutions written (default = number of available alternate solutions)</td>
</tr>
<tr>
<td>ams_stub</td>
<td>stub for alternate MIP solutions. The number of alternative MIP solution files written is determined by three keywords:  ams_limit gives the maximum number of files written;  ams_eps gives a relative tolerance on the objective values of alternative solutions;  ams_epsabs gives an absolute tolerance on how much worse the objectives can be.</td>
</tr>
</tbody>
</table>
Introductory Example

Multicommodity transportation . . .
- Products available at factories
- Products needed at stores
- Plan shipments at lowest cost

. . . with practical restrictions
- Cost has fixed and variable parts
- Shipments cannot be too small
- Factories cannot serve too many stores
Multicommodity Transportation

Given

\( O \) Set of origins (factories)
\( D \) Set of destinations (stores)
\( P \) Set of products

and

\( a_{ip} \) Amount available, for each \( i \in O \) and \( p \in P \)
\( b_{jp} \) Amount required, for each \( j \in D \) and \( p \in P \)
\( l_{ij} \) Limit on total shipments, for each \( i \in O \) and \( j \in D \)
\( c_{ijp} \) Shipping cost per unit, for each \( i \in O, j \in D, p \in P \)
\( d_{ij} \) Fixed cost for shipping any amount from \( i \in O \) to \( j \in D \)
\( s \) Minimum total size of any shipment
\( n \) Maximum number of destinations served by any origin
Multicommodity Transportation

Mathematical Formulation

Determine

\[ X_{ijp} \] Amount of each \( p \in P \) to be shipped from \( i \in O \) to \( j \in D \)
\[ Y_{ij} \] 1 if any product is shipped from \( i \in O \) to \( j \in D \)
0 otherwise

To minimize

\[ \sum_{i \in O} \sum_{j \in D} \sum_{p \in P} c_{ijp} X_{ijp} + \sum_{i \in O} \sum_{j \in D} d_{ij} Y_{ij} \]

Total variable cost plus total fixed cost
Multicommodity Transportation

Mathematical Formulation

Subject to

\[ \sum_{j \in D} X_{ijp} \leq a_{ip} \quad \text{for all } i \in O, p \in P \]

Total shipments of product \( p \) out of origin \( i \)

must not exceed availability

\[ \sum_{i \in O} X_{ijp} = b_{jp} \quad \text{for all } j \in D, p \in P \]

Total shipments of product \( p \) into destination \( j \)

must satisfy requirements
Mathematical Formulation

\textit{Subject to}

1. \[ \sum_{p \in P} X_{ijp} \leq l_{ij} Y_{ij} \quad \text{for all } i \in O, j \in D \]
   
   When there are shipments from origin \( i \) to destination \( j \),
   the total may not exceed the limit, and \( Y_{ij} \) must be 1

2. \[ \sum_{p \in P} X_{ijp} \geq s Y_{ij} \quad \text{for all } i \in O, j \in D \]
   
   When there are shipments from origin \( i \) to destination \( j \),
   the total amount of shipments must be at least \( s \)

3. \[ \sum_{j \in D} Y_{ij} \leq n \quad \text{for all } i \in O \]
   
   Number of destinations served by origin \( i \)
   must be as most \( n \)
Multicommodity Transportation

AMPL Formulation

Symbolic data

<table>
<thead>
<tr>
<th>Command</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>set ORIG;</td>
<td># origins</td>
</tr>
<tr>
<td>set DEST;</td>
<td># destinations</td>
</tr>
<tr>
<td>set PROD;</td>
<td># products</td>
</tr>
<tr>
<td>param supply {ORIG,PROD} &gt;= 0</td>
<td># availabilities at origins</td>
</tr>
<tr>
<td>param demand {DEST,PROD} &gt;= 0</td>
<td># requirements at destinations</td>
</tr>
<tr>
<td>param limit {ORIG,DEST} &gt;= 0</td>
<td># capacities of links</td>
</tr>
<tr>
<td>param vcost {ORIG,DEST,PROD}</td>
<td># variable shipment cost</td>
</tr>
<tr>
<td>param fcost {ORIG,DEST} &gt; 0</td>
<td># fixed usage cost</td>
</tr>
<tr>
<td>param minload &gt;= 0</td>
<td># minimum shipment size</td>
</tr>
<tr>
<td>param maxserve integer &gt; 0</td>
<td># maximum destinations served</td>
</tr>
</tbody>
</table>
Multicommodity Transportation

AMPL Formulation

Symbolic model: variables and objective

\[
\text{var } \textbf{Trans} \{\text{ORIG,DEST,PROD}\} \geq 0; \quad \# \text{ actual units to be shipped}
\]

\[
\text{var } \textbf{Use} \{\text{ORIG, DEST}\} \text{ binary}; \quad \# 1 \text{ if link used, 0 otherwise}
\]

\[
\text{minimize } \textbf{Total\_Cost}:
\]

\[
\sum \{i \text{ in ORIG, } j \text{ in DEST, } p \text{ in PROD}\} \text{vcost}[i,j,p] \times \text{Trans}[i,j,p]
\]

\[
+ \sum \{i \text{ in ORIG, } j \text{ in DEST}\} \text{fcost}[i,j] \times \text{Use}[i,j];
\]

\[
\sum_{i \in O} \sum_{j \in D} \sum_{p \in P} c_{ijp} X_{ijp} + \sum_{i \in O} \sum_{j \in D} d_{ij} Y_{ij}
\]
Multicommodity Transportation

AMPL Formulation

Symbolic model: constraint

subject to Supply \{i in ORIG, p in PROD\}:
    sum \{j in DEST\} Trans[i,j,p] <= supply[i,p];

\[ \sum_{j \in D} X_{ijp} \leq a_{ip}, \text{ for all } i \in O, p \in P \]
Multicommodity Transportation

AMPL Formulation

Symbolic model: constraints

subject to Supply {i in ORIG, p in PROD}:
    sum {j in DEST} Trans[i,j,p] <= supply[i,p];

subject to Demand {j in DEST, p in PROD}:
    sum {i in ORIG} Trans[i,j,p] = demand[j,p];

subject to Multi {i in ORIG, j in DEST}:
    sum {p in PROD} Trans[i,j,p] <= limit[i,j] * Use[i,j];

subject to Min_Ship {i in ORIG, j in DEST}:
    sum {p in PROD} Trans[i,j,p] >= minload * Use[i,j];

subject to Max_Serve {i in ORIG}:
    sum {j in DEST} Use[i,j] <= maxserve;
Multicommodity Transportation

AMPL Formulation

Explicit data independent of symbolic model

set ORIG := GARY CLEV PITT ;
set DEST := FRA DET LAN WIN STL FRE LAF ;
set PROD := bands coils plate ;
param supply (tr):  GARY   CLEV   PITT :=
   bands    400    700    800
   coils    800   1600   1800
   plate    200    300    300 ;

param demand (tr):
   FRA   DET   LAN   WIN   STL   FRE   LAF :=
   bands   300   300   100    75   650   225   250
   coils   500   750   400   250   950   850   500
   plate   100   100    0    50   200   100   250 ;

param limit default 625 ;
param minload := 375 ;
param maxserve := 5 ;
### AMPL Formulation

#### Explicit data (continued)

```AMPL
param vcost :=
  [*,*,bands]: FRA DET LAN WIN STL FRE LAF :=
   GARY  30  10   8  10  11  71  6
   CLEV  22  7  10   7  21  82 13
   PITT  19 11  12  10  25  83 15

  [*,*,coils]: FRA DET LAN WIN STL FRE LAF :=
   GARY  39 14  11  14  16  82  8
   CLEV  27  9  12   9  26  95 17
   PITT  24 14  17  13  28  99 20

  [*,*,plate]: FRA DET LAN WIN STL FRE LAF :=
   GARY  41 15  12  16  17  86  8
   CLEV  29  9  13   9  28  99 18
   PITT  26 14  17  13  31 104 20;

param fcost: FRA DET LAN WIN STL FRE LAF :=
  GARY  3000 1200 1200 1200 2500 3500 2500
  CLEV  2000 1000 1500 1200 2500 3000 2200
  PITT  2000 1200 1500 1500 2500 3500 2200;
```

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Multicommodity Transportation

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Multicommodity Transportation

AMPL Solution

Model + data = problem instance to be solved

```
AMPL: model multipG.mod;
AMPL: data multipG.dat;
AMPL: option solver gurobi;
AMPL: solve;
Gurobi 5.0.0: optimal solution; objective 235625
394 simplex iterations
46 branch-and-cut nodes
AMPL: display Use;
Use [*,*]:
   DET  FRA  FRE  LAF  LAN  STL  WIN  :=
  CLEV  1   1   1   0   1   1   0
  GARY  0   0   0   1   0   1   1
  PITT  1   1   1   1   0   1   0
```

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Multicommodity Transportation

AMPL Solution

Solver choice independent of model and data

```AMPL
ampl: model multipG.mod;
ampl: data multipG.dat;
ampl: option solver cplex;
ampl: solve;
Cplex 12.4.0.0: optimal integer solution; objective 235625
394 MIP simplex iterations
41 branch-and-bound nodes
ampl: display Use;
Use [*,*]
  :   DET FRA FRE LAF LAN STL WIN   :=
CLEV  1  1  1  0  1  1  0
GARY  0  0  0  1  0  1  1
PITT  1  1  1  1  0  1  0
;
```
Multicommodity Transportation

AMPL Solution

Examine results

```ampl
ampl: display {i in ORIG, j in DEST}

ampl: sum {p in PROD} Trans[i,j,p] / limit[i,j];

:      DET   FRA    FRE    LAF    LAN    STL    WIN    :=
CLEV   1      0.6    0.88   0      0.8    0.88   0      GARY   0      0      0      0.64   0     1      0.6
PITT   0.84   0.84   1      0.96   0     1      0      :=
                  0.96   0     1      0
WIN    0.85   0.85   1      0.95   0     1      0.6

ampl: display Max_Serve.body;

CLEV  5
GARY  3
PITT  5

ampl: display TotalCost,

ampl: sum {i in ORIG, j in DEST} fcost[i,j] * Use[i,j];

TotalCost = 235625
sum {i in ORIG, j in DEST} fcost[i,j] * Use[i,j] = 27600
```
Indexed over sets of pairs and triples

```plaintext
set ORIG;   # origins
set DEST;   # destinations
set PROD;   # products
set SHIP within {ORIG,DEST,PROD};
            # (i,j,p) in SHIP ==> can ship p from i to j
set LINK = setof {(i,j,p) in SHIP} (i,j);
            # (i,j) in LINK ==> can ship some products from i to j

var Trans {SHIP} >= 0;   # actual units to be shipped
var Use {LINK} binary;   # 1 if link used, 0 otherwise

minimize Total_Cost:
    sum {(i,j,p) in SHIP} vcost[i,j,p] * Trans[i,j,p]
+ sum {(i,j) in LINK} fcost[i,j] * Use[i,j];
```
Multicommodity Transportation

AMPL “Sparse” Network

Constraint for dense network

subject to Supply \{i \in ORIG, p \in PROD\}:

\[ \sum \{j \in DEST\} Trans[i,j,p] \leq supply[i,p]; \]

Constraint for sparse network

subject to Supply \{i \in ORIG, p \in PROD\}:

\[ \sum \{(i,j,p) \in SHIP\} Trans[i,j,p] \leq supply[i,p]; \]
Multicommodity Transportation

AMPL “Sparse” Network

All constraints

subject to Supply {i in ORIG, p in PROD}:
    \( \sum \{(i,j,p) \in \text{SHIP}\} \ Trans[i,j,p] \leq \text{supply}[i,p] \);

subject to Demand {j in DEST, p in PROD}:
    \( \sum \{(i,j,p) \in \text{SHIP}\} \ Trans[i,j,p] = \text{demand}[j,p] \);

subject to Multi {i in ORIG, j in DEST}:
    \( \sum \{(i,j,p) \in \text{SHIP}\} \ Trans[i,j,p] \leq \text{limit}[i,j] \times \text{Use}[i,j] \);

subject to Min_Ship {i in ORIG, j in DEST}:
    \( \sum \{(i,j,p) \in \text{SHIP}\} \ Trans[i,j,p] \geq \text{minload} \times \text{Use}[i,j] \);

subject to Max_Serve {i in ORIG}:
    \( \sum \{(i,j) \in \text{LINK}\} \ \text{Use}[i,j] \leq \text{maxserve} \);
Multicommodity Transportation

AMPL “Sparse” Network

1st dataset: shipments allowed

```
set SHIP :=
    (*,*,bands): FRA DET LAN WIN STL FRE LAF :=
    GARY    +  +  +  +  +  -  +    
    CLEV    +  -  +  -  +  +  +  
    PITT    -  +  +  +  +  +  +  
    (*,*,coils): FRA DET LAN WIN STL FRE LAF :=
    GARY    +  +  +  +  +  +  +  -    
    CLEV    +  +  -  +  +  +  +  +    
    PITT    +  +  +  +  +  +  +  +    
    (*,*,plate): FRA DET LAN WIN STL FRE LAF :=
    GARY    +  +  -  +  +  -  +  +    
    CLEV    +  +  +  +  +  +  +  +    
    PITT    -  +  +  -  +  +  +  +    
```

Multicommodity Transportation
**Multicommodity Transportation**

**AMPL “Sparse” Network**

2nd dataset: shipments allowed

```plaintext
set SHIP :=
    (*,*,bands):  FRA  DET  LAN  WIN  STL  FRE  LAF :=
                        GARY    +    +    +    +    +    -    -
                        CLEV    -    +    +    -    +    +    +
                        PITT    +    -    +    +    +    +    +
    (*,*,coils):  FRA  DET  LAN  WIN  STL  FRE  LAF :=
                        GARY    +    +    +    +    +    +    +
                        CLEV    +    +    -    +    +    +    +
                        PITT    +    +    +    +    +    +    +
    (*,*,plate):  FRA  DET  LAN  WIN  STL  FRE  LAF :=
                        GARY    -    +    +    +    +    -    +
                        CLEV    +    +    +    +    +    +    +
                        PITT    +    +    -    -    +    +    +
```

Multicommodity Transportation

AMPL “Sparse” Network

Same model, different data

```
AMPL: model multmipT.mod;
AMPL: data multmipT1.dat;
AMPL: solve;
Gurobi 4.6.0: optimal solution; objective 247725
108 simplex iterations
13 branch-and-cut nodes
AMPL: reset data;
AMPL: data multmipT2.dat;
AMPL: solve;
Gurobi 4.6.0: optimal solution; objective 237775
79 simplex iterations
AMPL:
```
1: Parametric Analysis

Try different limits on destinations served

- Reduce parameter maxserve and re-solve
  * until there is no feasible solution
- Display results
  * parameter value
  * numbers of destinations actually served

Try different supplies of plate at Gary

- Increase parameter supply['GARY','plate'] and re-solve
  * until dual is zero (constraint is slack)
- Record results
  * distinct dual values
  * corresponding objective values

... display results at the end
Parametric Analysis on limits

Script to test sensitivity to serve limit

model multmipG.mod;
data multmipG.dat;

option solver gurobi;
for {m in 7..1 by -1} {
  let maxserve := m;
solve;
  if solve_result = 'infeasible' then break;
  display maxserve, Max_Serve.body;
}
Parametric Analysis on limits

Run showing sensitivity to serve limit

```ampl
ampl: include multmipServ.run;

Gurobi 4.6.0: optimal solution; objective 233150
maxserve = 7
CLEV 5  GARY 3  PITT 6

Gurobi 4.6.0: optimal solution; objective 233150
maxserve = 6
CLEV 5  GARY 3  PITT 6

Gurobi 4.6.0: optimal solution; objective 235625
maxserve = 5
CLEV 5  GARY 3  PITT 5

Gurobi 4.6.0: infeasible
```
Parametric Analysis on supplies

Script to test sensitivity to plate supply at GARY

```
set SUPPLY default {};  
param sup_obj {SUPPLY};  
param sup_dual {SUPPLY};  
let supply['GARY','plate'] := 200;  
param sup_step = 10;  
param previous_dual default -Infinity;  
repeat while previous_dual < 0 {
    solve;  
    if Supply['GARY','plate'].dual > previous_dual then {
        let SUPPLY := SUPPLY union {supply['GARY','plate']};  
        let sup_obj[supply['GARY','plate']] := Total_Cost;  
        let sup_dual[supply['GARY','plate']] := Supply['GARY','plate'].dual;  
        let previous_dual := Supply['GARY','plate'].dual;  
    }
    let supply['GARY','plate'] := supply['GARY','plate'] + supply_step;
}
```
Parametric Analysis on supplies

Run showing sensitivity to plate supply at GARY

ampl: include mult mipSupply.run;

ampl: display sup_obj, sup_dual;

<table>
<thead>
<tr>
<th>sup_obj</th>
<th>sup_dual</th>
</tr>
</thead>
<tbody>
<tr>
<td>200</td>
<td>223504</td>
</tr>
<tr>
<td>380</td>
<td>221171</td>
</tr>
<tr>
<td>460</td>
<td>220260</td>
</tr>
<tr>
<td>510</td>
<td>219754</td>
</tr>
<tr>
<td>560</td>
<td>219413</td>
</tr>
</tbody>
</table>

;
Parametric: Observations

Results of solve can be tested

- Check whether problem is no longer feasible
  - if solve_result = 'infeasible' then break;

Parameters are true objects

- Assign new value to param supply
  - let supply['GARY','plate'] :=
    supply['GARY','plate'] + supply_step;
- Problem instance changes accordingly

Sets are true data

- Assign new value to set SUPPLY
  - let SUPPLY := SUPPLY union {supply['GARY','plate']};
- All indexed entities change accordingly
2a: Solution Generation \textit{via Cuts}

\textit{Same multicommodity transportation model}

\textit{Generate n best solutions using different routes}

- Display routes used by each solution
Solutions via Cuts

Script

```plaintext
param nSols default 0;
param maxSols = 3;
model multimipG.mod;
data multimipG.dat;
set USED {1..nSols} within {ORIG,DEST};
subject to exclude {k in 1..nSols}:
    sum {(i,j) in USED[k]} (1-Use[i,j]) +
    sum {(i,j) in {ORIG,DEST} diff USED[k]} Use[i,j] >= 1;
repeat {
    solve;
    display Use;
    let nSols := nSols + 1;
    let USED[nSols] := {i in ORIG, j in DEST: Use[i,j] > .5};
} until nSols = maxSols;
```
AMPL Scripting

Run showing 3 best solutions

```ampl
ampl: include multmipBestA.run;

Gurobi 4.6.0: optimal solution; objective 235625
:  DET  FRA  FRE  LAF  LAN  STL  WIN  :=
  CLEV  1  1  0  1  1  0
  GARY  0  0  1  0  1  1
  PITT  1  1  1  0  1  0;

Gurobi 4.6.0: optimal solution; objective 237125
:  DET  FRA  FRE  LAF  LAN  STL  WIN  :=
  CLEV  1  1  1  0  1  0
  GARY  0  1  0  1  0  1
  PITT  1  1  1  0  1  0;

Gurobi 4.6.0: optimal solution; objective 238225
:  DET  FRA  FRE  LAF  LAN  STL  WIN  :=
  CLEV  1  0  1  0  1  1
  GARY  0  1  0  1  0  1
  PITT  1  1  1  0  1  0;
```
Solutions via Cuts: Observations

Same expressions describe sets and indexing

- Index a summation
  
  ```
  ... sum {(i,j) in {ORIG,DEST} diff USED[k]} Use[i,j] >= 1;
  ```

- Assign a value to a set
  
  ```
  let USED[nSols] := {i in ORIG, j in DEST: Use[i,j] > .5};
  ```

New cuts defined automatically

- Index cuts over a set
  
  ```
  subject to exclude {k in 1..nSols}: ...
  ```

- Add a cut by expanding the set
  
  ```
  let nSols := nSols + 1;
  ```
2b: Solution Generation via Solver

Same model

Ask solver to return multiple solutions

- Set options
- Get all results from one “solve”
- Retrieve and display each solution
Solutions via Solver

Script

```plaintext
option solver cplex;
option cplex_options "poolstub=multmip poolcapacity=3 \n    populate=1 poolintensity=4 poolreplace=1";

solve;

for {i in 1..Current.npool} {
    solution ("multmip" & i & ".sol");
    display Use;
}
```
Solutions *via Solver*

**Results**

```ampl
ampl: include multimipBestB.run;

CPLEX 12.4.0.0: poolstub=multmip
poolcapacity=3
populate=1
poolintensity=4
poolreplace=1

CPLEX 12.4.0.0: optimal integer solution; objective 235625
439 MIP simplex iterations
40 branch-and-bound nodes

Wrote 3 solutions in solution pool
to files multimip1.sol ... multimip3.sol.

Suffix npool OUT;
```
Solutions via Solver

Results (continued)

Solution pool member 1 (of 3); objective 235625

: DET FRA FRE LAF LAN STL WIN :=
CLEV  1  1  1  0  1  1  0
GARY  0  0  0  1  0  1  1
PITT  1  1  1  1  0  1  0;

Solution pool member 2 (of 3); objective 238225

: DET FRA FRE LAF LAN STL WIN :=
CLEV  1  0  1  0  1  1  1
GARY  0  1  0  1  0  1  0
PITT  1  1  1  1  0  1  0;

Solution pool member 3 (of 3); objective 237125

: DET FRA FRE LAF LAN STL WIN :=
CLEV  1  1  1  1  0  1  0
GARY  0  0  0  1  0  1  1
PITT  1  1  1  0  1  1  0;
Solutions via Solver: Observations

Filenames can be formed dynamically

- Write a (string expression)
- Numbers are automatically converted
  * solution ("multmip" & i & ".sol");
3: Heuristic Optimization

Workforce planning

- Cover demands for workers
  - Each “shift” requires a certain number of employees
  - Each employee works a certain “schedule” of shifts

- Satisfy scheduling rules
  - Only “valid” schedules from given list may be used
  - Each schedule that is used at all must be worked by at least ?? employees

- Minimize total workers needed
  - Which schedules should be used?
  - How many employees should work each schedule?

Difficult instances

- Set ?? to a “hard” value
- Get a very good solution quickly
Heuristic

Model (sets, parameters)

```
set SHIFTS;               # shifts
param Nsched;             # number of schedules;
set SCHEDS = 1..Nsched;   # set of schedules
set SHIFT_LIST {SCHEDS} within SHIFTS;

param rate {SCHEDS} >= 0;       # pay rates
param required {SHIFTS} >= 0;   # staffing requirements
param least_assign >= 0;        # min workers on any schedule used
```
Heuristic

Model (variables, objective, constraints)

var Work \{SCHEDS\} \geq 0 \text{ integer};
var Use \{SCHEDS\} \geq 0 \text{ binary};

minimize Total_Cost:
\quad \text{sum } \{j \text{ in SCHEDS}\} \text{ rate}[j] * \text{Work}[j];

subject to Shift_Needs \{i \text{ in SHIFTS}\}:
\quad \text{sum } \{j \text{ in SCHEDS: } i \text{ in SHIFT_LIST}[j]\} \text{ Work}[j] \geq \text{required}[i];

subject to Least_Use1 \{j \text{ in SCHEDS}\}:
\quad \text{least_assign } * \text{Use}[j] \leq \text{Work}[j];

subject to Least_Use2 \{j \text{ in SCHEDS}\}:
\quad \text{Work}[j] \leq (\max \{i \text{ in SHIFT_LIST}[j]\} \text{ required}[i]) * \text{Use}[j];
Heuristic

Data

set SHIFTS := Mon1 Tue1 Wed1 Thu1 Fri1 Sat1
            Mon2 Tue2 Wed2 Thu2 Fri2 Sat2
            Mon3 Tue3 Wed3 Thu3 Fri3 ;

param Nsched := 126 ;

set SHIFT_LIST[1] := Mon1 Tue1 Wed1 Thu1 Fri1 ;
set SHIFT_LIST[2] := Mon1 Tue1 Wed1 Thu1 Fri2 ;
set SHIFT_LIST[3] := Mon1 Tue1 Wed1 Thu1 Fri3 ;
set SHIFT_LIST[4] := Mon1 Tue1 Wed1 Thu1 Sat1 ;
set SHIFT_LIST[5] := Mon1 Tue1 Wed1 Thu1 Sat2 ; .......

param required :=
            Mon1 100  Mon2  78  Mon3  52
            Tue1 100  Tue2  78  Tue3  52
            Wed1 100  Wed2  78  Wed3  52
            Thu1 100  Thu2  78  Thu3  52
            Fri1 100  Fri2  78  Fri3  52
            Sat1 100  Sat2  78 ;
Heuristic

**Hard case:** `least_assign = 19`

```ampl
ampl: model sched1.mod;
ampl: data sched.dat;
ampl: let least_assign := 19;
ampl: option solver cplex;
ampl: solve;

CPLEX 12.2.0.2: optimal integer solution; objective 269
635574195 MIP simplex iterations
86400919 branch-and-bound nodes

ampl: option omit_zero_rows 1, display_1col 0;
ampl: display Work;

Work [*] :=
    4 22    16 39    55 39    78 39   101 39   106 52   122 39   ... 94.8 minutes
```

Robert Fourer, Alternatives for Scripting in an Algebraic Modeling Language
EURO Vilnius — 8-11 July, 2012 — WD33 Optimization Modeling II
Heuristic

*Alternative, indirect approach*

- Step 1: Relax integrality of *Work* variables
  Solve for zero-one *Use* variables

- Step 2: Fix *Use* variables
  Solve for integer *Work* variables

  . . . *not necessarily optimal, but . . .*
Heuristic

Script

```plaintext
model sched1.mod;
data sched.dat;
let least_assign := 19;

let {j in SCHEDS} Work[j].relax := 1;
solve;

fix {j in SCHEDS} Use[j];
let {j in SCHEDS} Work[j].relax := 0;
solve;
```
Heuristic

Results

ampl: include sched1-fix.run;

CPLEX 12.2.0.2: optimal integer solution; objective 268.5
32630436 MIP simplex iterations
2199508 branch-and-bound nodes

Work [*] :=
  1 24  32 19  80 19.5  107 33  126 19.5
  3 19  66 19  90 19.5  109 19
  10 19  72 19.5  105 19.5  121 19;

CPLEX 12.2.0.2: optimal integer solution; objective 269
2 MIP simplex iterations
0 branch-and-bound nodes

Work [*] :=
  1 24  10 19  66 19  80 19  105 20  109 19  126 20
  3 19  32 19  72 19  90 20  107 33  121 19;

... 2.85 minutes
Heuristic: Observations

Models can be changed dynamically

- Adapt modeling expressions
- Execute model-related commands
  - fix {j in SCHEDS} Use[j];
- Assign values to properties of model components
  - let {j in SCHEDS} Work[j].relax := 1;
4: Pattern Generation

Roll cutting
- Min rolls cut (or material wasted)
- Decide number of each pattern to cut
- Meet demands for each ordered width

Generate cutting patterns
- Read general model
- Read data: demands, raw width
- Compute data: all usable patterns
- Solve problem instance
Pattern Generation

Model

```plaintext
param roll_width > 0;
set WIDTHS ordered by reversed Reals;
param orders {WIDTHS} > 0;
param maxPAT integer >= 0;
param nPAT integer >= 0, <= maxPAT;
param nbr {WIDTHS,1..maxPAT} integer >= 0;

var Cut {1..nPAT} integer >= 0;

minimize Number:
    sum {j in 1..nPAT} Cut[j];

subj to Fulfill {i in WIDTHS}:
    sum {j in 1..nPAT} nbr[i,j] * Cut[j] >= orders[i];
```
Pattern Generation

Data

<table>
<thead>
<tr>
<th>Width</th>
<th>Order</th>
</tr>
</thead>
<tbody>
<tr>
<td>60</td>
<td>3</td>
</tr>
<tr>
<td>30</td>
<td>21</td>
</tr>
<tr>
<td>25.5</td>
<td>94</td>
</tr>
<tr>
<td>20</td>
<td>50</td>
</tr>
<tr>
<td>17.25</td>
<td>288</td>
</tr>
<tr>
<td>15</td>
<td>178</td>
</tr>
<tr>
<td>12.75</td>
<td>112</td>
</tr>
<tr>
<td>10</td>
<td>144</td>
</tr>
</tbody>
</table>
Pattern Generation

Script (initialize)

```plaintext
model cutPAT.mod;
data ChvatalD.dat;

model;
param curr_sum >= 0;
param curr_width > 0;
param pattern {WIDTHS} integer >= 0;
let maxPAT := 100000000;
let nPAT := 0;
let curr_sum := 0;
let curr_width := first(WIDTHS);
let {w in WIDTHS} pattern[w] := 0;
```
Pattern Generation

Script (loop)

repeat {
    if curr_sum + curr_width <= roll_width then {
        let pattern[curr_width] := floor((roll_width-curr_sum)/curr_width);
        let curr_sum := curr_sum + pattern[curr_width] * curr_width;
    }
    if curr_width != last(WIDTHS) then
        let curr_width := next(curr_width,WIDTHS);
    else {
        let nPAT := nPAT + 1;
        let {w in WIDTHS} nbr[w,nPAT] := pattern[w];
        let curr_sum := curr_sum - pattern[last(WIDTHS)] * last(WIDTHS);
        let pattern[last(WIDTHS)] := 0;
        let curr_width := min {w in WIDTHS: pattern[w] > 0} w;
        if curr_width < Infinity then {
            let curr_sum := curr_sum - curr_width;
            let pattern[curr_width] := pattern[curr_width] - 1;
            let curr_width := next(curr_width,WIDTHS);
        }
        else break;
    }
}
Pattern Generation

Script (solve, report)

```c
option solver gurobi;
solve;
printf "\n%5i patterns, %3i rolls", nPAT, sum {j in 1..nPAT} Cut[j];
printf "\n\n Cut  ";
printf {j in 1..nPAT: Cut[j] > 0}: "%3i", Cut[j];
printf "\n\n";
for {i in WIDTHS} {
    printf "%7.2f ", i;
    printf {j in 1..nPAT: Cut[j] > 0}: "%3i", nbr[i,j];
    printf "\n";
}
printf "\nWASTE = %5.2f\n\n",
    100 * (1 - (sum {i in WIDTHS} i * orders[i]) / (roll_width * Number));
```
Pattern Generation

Results

ampl: include cutPatEnum.run

Gurobi 4.6.1: optimal solution; objective 164
15 simplex iterations

290 patterns, 164 rolls

<table>
<thead>
<tr>
<th>Cut</th>
<th>3</th>
<th>7</th>
<th>50</th>
<th>44</th>
<th>17</th>
<th>25</th>
<th>2</th>
<th>16</th>
</tr>
</thead>
<tbody>
<tr>
<td>60.00</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>30.00</td>
<td>0</td>
<td>3</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>25.50</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>20.00</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>3</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>17.25</td>
<td>0</td>
<td>0</td>
<td>3</td>
<td>2</td>
<td>0</td>
<td>2</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>15.00</td>
<td>2</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>2</td>
<td>2</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>12.75</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>7</td>
<td>0</td>
</tr>
<tr>
<td>10.00</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>9</td>
</tr>
</tbody>
</table>

WASTE = 0.32%
Pattern Generation

Data 2

```
param roll_width := 349;
param: WIDTHS: orders :=
    28.75  7
    33.75 23
    34.75 23
    37.75 31
    38.75 10
    39.75 39
    40.75 58
    41.75 47
    42.25 19
    44.75 13
    45.75 26;
```
Pattern Generation

Results 2

ampl: include cutPatEnum.run

Gurobi 4.6.1: optimal solution; objective 34
291 simplex iterations

54508 patterns, 34 rolls

<table>
<thead>
<tr>
<th>Cut</th>
<th>8</th>
<th>1</th>
<th>1</th>
<th>1</th>
<th>3</th>
<th>1</th>
<th>1</th>
<th>1</th>
<th>1</th>
<th>2</th>
<th>7</th>
<th>2</th>
<th>3</th>
<th>1</th>
<th>1</th>
</tr>
</thead>
<tbody>
<tr>
<td>45.75</td>
<td>3</td>
<td>2</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>44.75</td>
<td>1</td>
<td>2</td>
<td>2</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>42.25</td>
<td>0</td>
<td>2</td>
<td>0</td>
<td>0</td>
<td>4</td>
<td>2</td>
<td>2</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>41.75</td>
<td>4</td>
<td>2</td>
<td>0</td>
<td>2</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>1</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>40.75</td>
<td>0</td>
<td>0</td>
<td>4</td>
<td>4</td>
<td>1</td>
<td>4</td>
<td>3</td>
<td>0</td>
<td>2</td>
<td>3</td>
<td>1</td>
<td>6</td>
<td>3</td>
<td>2</td>
<td>2</td>
</tr>
<tr>
<td>39.75</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>0</td>
<td>0</td>
<td>5</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>0</td>
</tr>
<tr>
<td>38.75</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>4</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>3</td>
</tr>
<tr>
<td>37.75</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>4</td>
<td>0</td>
<td>0</td>
<td>6</td>
<td>2</td>
<td>4</td>
</tr>
<tr>
<td>34.75</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>4</td>
<td>0</td>
<td>3</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>3</td>
<td>0</td>
<td>1</td>
<td>0</td>
</tr>
<tr>
<td>33.75</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>3</td>
<td>0</td>
<td>4</td>
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<td>1</td>
<td>2</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
<tr>
<td>28.75</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>2</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>2</td>
<td>1</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
<td>0</td>
</tr>
</tbody>
</table>

WASTE = 0.69%
Pattern Generation

Data 3

```
param roll_width := 172 ;
param: WIDTHS: orders :=
  25.000   5
  24.750   73
  18.000   14
  17.500   4
  15.500   23
  15.375   5
  13.875   29
  12.500   87
  12.250   9
  12.000   31
  10.250   6
  10.125   14
  10.000   43
  8.750    15
  8.500    21
  7.750    5 ;
```
Pattern Generation

Results 3 (using a subset of patterns)

<table>
<thead>
<tr>
<th>Cut</th>
<th>1 1 1 1 4 4 4 1 1 2 5 2 1 1 1 3</th>
</tr>
</thead>
<tbody>
<tr>
<td>25.00</td>
<td>2 1 1 1 0 0 0 0 0 0 0 0 0 0 0 0</td>
</tr>
<tr>
<td>24.75</td>
<td>1 2 1 0 5 4 3 2 2 2 2 1 1 0 0 0</td>
</tr>
<tr>
<td>18.00</td>
<td>0 0 0 0 1 0 0 1 0 0 0 1 1 5 1 0</td>
</tr>
<tr>
<td>17.50</td>
<td>0 3 0 0 0 0 0 0 0 0 0 0 0 0 1 0</td>
</tr>
<tr>
<td>.......</td>
<td></td>
</tr>
<tr>
<td>10.12</td>
<td>0 2 0 0 0 1 2 0 0 0 0 0 0 0 0 0</td>
</tr>
<tr>
<td>10.00</td>
<td>0 0 0 0 0 0 2 0 1 3 0 6 0 0 2 0</td>
</tr>
<tr>
<td>8.75</td>
<td>0 0 1 0 0 0 0 0 0 2 0 2 0 0 0 2</td>
</tr>
<tr>
<td>8.50</td>
<td>0 0 2 0 0 2 0 0 0 0 0 4 3 0 0 0</td>
</tr>
<tr>
<td>7.75</td>
<td>0 0 0 0 1 0 0 1 0 0 0 0 0 0 0 0</td>
</tr>
</tbody>
</table>

WASTE = 0.62%
Pattern Generation: Observations

Parameters can serve as script variables

- Declare as in model
  
  ```plaintext
  * param pattern {WIDTHS} integer >= 0;
  ```

- Use in algorithm
  
  ```plaintext
  * let pattern[curr_width] := pattern[curr_width] - 1;
  ```

- Assign to model parameters
  
  ```plaintext
  * let {w in WIDTHS} nbr[w,nPAT] := pattern[w];
  ```

Scripts are easy to modify

- Store only every 100th pattern found
  
  ```plaintext
  * if nPAT mod 100 = 0 then
    let {w in WIDTHS} nbr[w,nPAT/100] := pattern[w];
  ```
5: Decomposition

Stochastic nonlinear location-transportation

- Min expected total cost
  - Nonlinear construction costs at origins
  - Linear transportation costs from origins to destinations
- Stochastic demands with recourse
  - Decide what to build
  - Observe demands and decide what to ship

Solve by Benders decomposition

- Nonlinear master problem
- Linear subproblem for each scenario
Decomposition

Original model (sets, parameters, variables)

```plaintext
set WHSE;   # shipment origins (warehouses)
set STOR;   # shipment destinations (stores)

param build_cost {WHSE} > 0;  # costs per unit to build warehouse
param build_limit {WHSE} > 0; # limits on units shipped

var Build {i in WHSE} >= 0, <= .9999 * build_limit[i];
   # capacities of warehouses to build

set SCEN;                    # demand scenarios

param prob {SCEN} >= 0, <= 1;   # probabilities of scenarios
param demand {STOR,SCEN} >= 0;  # amounts required at stores

param ship_cost {WHSE,STOR} >= 0; # shipment costs per unit

var Ship {WHSE,STOR,SCEN} >= 0;  # amounts to be shipped
```
Decomposition

Original model (objective, constraints)

\[
\text{minimize Total\_Cost:}
\]
\[
\quad \text{sum \{i in WHSE\}}
\]
\[
\quad \quad \text{build\_cost[i] * Build[i] / (1 - Build[i]/build\_limit[i]) +}
\]
\[
\quad \quad \text{sum \{s in SCEN\} prob[s] *}
\]
\[
\quad \quad \quad \text{sum \{i in WHSE, j in STOR\} ship\_cost[i,j] * Ship[i,j,s];}
\]

subj to Supply \{i in WHSE, s in SCEN\}:
\[
\quad \text{sum \{j in STOR\} Ship[i,j,s] <= Build[i];}
\]

subj to Demand \{j in STOR, s in SCEN\}:
\[
\quad \text{sum \{i in WHSE\} Ship[i,j,s] = demand[j,s];}
\]
Decomposition

Sub model (sets, parameters, variables)

set WHSE;       # shipment origins (warehouses)
set STOR;        # shipment destinations (stores)

param build {i in WHSE} >= 0, <= .9999 * build_limit[i];
                   # capacities of warehouses built

set SCEN;        # demand scenarios

param prob {SCEN} >= 0, <= 1;     # probabilities of scenarios
param demand {STOR,SCEN} >= 0;    # amounts required at stores

param ship_cost {WHSE,STOR} >= 0; # shipment costs per unit

var Ship {WHSE,STOR,SCEN} >= 0;    # amounts to be shipped
Decomposition

Sub model (objective, constraints)

```
param S symbolic in SCEN;

minimize Scen_Ship_Cost:
    prob[S] * sum {i in WHSE, j in STOR} ship_cost[i,j] * Ship[i,j];

subj to Supply {i in WHSE}:
    sum {j in STOR} Ship[i,j] <= build[i];

subj to Demand {j in STOR}:
    sum {i in WHSE} Ship[i,j] = demand[j,S];
```
Decomposition

Master model (sets, parameters, variables)

```
param build_cost {WHSE} > 0; # costs per unit to build warehouse
param build_limit {WHSE} > 0; # limits on units shipped
var Build {i in WHSE} >= 0, <= .9999 * build_limit[i]; # capacities of warehouses to build

param nCUT >= 0 integer;
param cut_type {SCEN,1..nCUT} symbolic
    within {"feas","infeas","none"};
param supply_price {WHSE,SCEN,1..nCUT} <= 0.000001;
param demand_price {STOR,SCEN,1..nCUT};

var Max_Exp_Ship_Cost {SCEN} >= 0;
```
Decomposition

Master model (objective, constraints)

\[
\text{minimize } \text{Expected_Total_Cost}:
\]
\[
\text{sum } \{i \text{ in WHSE}\} \text{ build_cost}[i] \times \text{Build}[i] / (1 - \text{Build}[i]/\text{build_limit}[i]) + \\
\text{sum } \{s \text{ in SCEN}\} \text{ Max_Exp_Ship_Cost}[s];
\]

subj to Cut_Defn \{s \text{ in SCEN, } k \text{ in 1..nCUT: cut_type}[s,k] != "none"\}:
\[
\text{if cut_type}[s,k] = "\text{feas}\" \text{ then } \text{Max_Exp_Ship_Cost}[s] \text{ else } 0 \geq \\
\text{sum } \{i \text{ in WHSE}\} \text{ supply_price}[i,s,k] \times \text{Build}[i] + \\
\text{sum } \{j \text{ in STOR}\} \text{ demand_price}[j,s,k] \times \text{demand}[j,s];
\]
Decomposition

Script (initialization)

model stbenders.mod;
data stnltrnloc.dat;
suffix dunbdd;
option presolve 0;

problem Sub: Ship, Scen_Ship_Cost, Supply, Demand;
  option solver cplex;
  option cplex_options 'primal presolve 0';
problem Master: Build, Max_Exp_Ship_Cost, Exp_Total_Cost, Cut_Defn;
  option solver minos;

let nCUT := 0;
param GAP default Infinity;
param RELGAP default Infinity;
param Exp_Ship_Cost;
Decomposition

Script (iteration)

repeat {
    solve Master;
    let \{i in WHSE\} build[i] := Build[i];
    let Exp_Ship_Cost := 0;
    let nCUT := nCUT + 1;
    for \{s in SCEN\} {
        let S := s;
        solve Sub;
        \textit{... generate a cut ...}
    }
    if forall \{s in SCEN\} cut_type[s,nCUT] != "infeas" then {
        let GAP := min (GAP,
            Exp_Ship_Cost - sum \{s in SCEN\} Max_Exp_Ship_Cost[s]);
        let RELGAP := 100 * GAP / Expected_Total_Cost;
    }
} until RELGAP <= .000001;
Decomposition

Script (cut generation)

```plaintext
for {s in SCEN} {
    let S := s;
    solve Sub;

    if Sub.result = "solved" then {
        let Exp_Ship_Cost := Exp_Ship_Cost + Scen_Ship_Cost;
        if Scen_Ship_Cost > Max_Exp_Ship_Cost[s] + 0.00001 then {
            let cut_type[s,nCUT] := "feas";
            let {i in WHSE} supply_price[i,s,nCUT] := Supply[i].dual;
            let {j in STOR} demand_price[j,s,nCUT] := Demand[j].dual;
        }
        else let cut_type[s,nCUT] := "none";
    }

    else if Sub.result = "infeasible" then {
        let cut_type[s,nCUT] := "infeas";
        let {i in WHSE} supply_price[i,s,nCUT] := Supply[i].dunbdd;
        let {j in STOR} demand_price[j,s,nCUT] := Demand[j].dunbdd;
    }
}
```
Decomposition

Results

ampl: include stbenders.run;

MASTER PROBLEM 1: 0.000000
SUB-PROBLEM 1 low: infeasible
SUB-PROBLEM 1 mid: infeasible
SUB-PROBLEM 1 high: infeasible
MASTER PROBLEM 2: 267806.267806
SUB-PROBLEM 2 low: 1235839.514234
SUB-PROBLEM 2 mid: 1030969.048921
SUB-PROBLEM 2 high: infeasible
MASTER PROBLEM 3: 718918.236014
SUB-PROBLEM 3 low: 1019699.661119
SUB-PROBLEM 3 mid: 802846.293052
SUB-PROBLEM 3 high: 695402.974379
GAP = 2517948.928551, RELGAP = 350.241349%
## Decomposition

### Results (continued)

<table>
<thead>
<tr>
<th>Master Problem</th>
<th>Lower Bound</th>
<th>Mid Point</th>
<th>Upper Bound</th>
</tr>
</thead>
<tbody>
<tr>
<td>4</td>
<td>2606868.719958</td>
<td></td>
<td></td>
</tr>
<tr>
<td>5</td>
<td>2685773.838398</td>
<td></td>
<td></td>
</tr>
<tr>
<td>6</td>
<td>2743483.001029</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

| Problem 4  | Low: 1044931.784272 | Mid: 885980.640150 | High: 944581.118758 |
| Problem 5  | Low: 1028785.052062 | Mid: 815428.531237 | High: 753627.189086 |
| Problem 6  | Low: 1000336.408156 | Mid: 785602.983289 | High: 725635.817601 |

GAP = 749765.716399, RELGAP = 28.761161%

GAP = 394642.837091, RELGAP = 14.693822%

GAP = 222288.965560, RELGAP = 8.102436%
Decomposition

Results (continued)

<table>
<thead>
<tr>
<th>Master Problem 7: 2776187.713412</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sub-Problem 7 low: 986337.500000</td>
</tr>
<tr>
<td>Sub-Problem 7 mid: 777708.466300</td>
</tr>
<tr>
<td>Sub-Problem 7 high: 693342.659287</td>
</tr>
<tr>
<td>GAP = 59240.084058, RELGAP = 2.133864%</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Master Problem 8: 2799319.395374</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sub-Problem 8 low: 991426.284976</td>
</tr>
<tr>
<td>Sub-Problem 8 mid: 777146.351060</td>
</tr>
<tr>
<td>Sub-Problem 8 high: 704353.854398</td>
</tr>
<tr>
<td>GAP = 38198.286498, RELGAP = 1.364556%</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Master Problem 9: 2814772.778136</th>
</tr>
</thead>
<tbody>
<tr>
<td>Sub-Problem 9 low: 987556.309573</td>
</tr>
<tr>
<td>Sub-Problem 9 mid: 772147.258329</td>
</tr>
<tr>
<td>Sub-Problem 9 high: 696060.666966</td>
</tr>
<tr>
<td>GAP = 17658.226624, RELGAP = 0.627341%</td>
</tr>
</tbody>
</table>
Decomposition

Results (continued)

<table>
<thead>
<tr>
<th>MASTER PROBLEM 10: 2818991.649514</th>
</tr>
</thead>
<tbody>
<tr>
<td>SUB-PROBLEM 10 mid: 771853.500000</td>
</tr>
<tr>
<td>SUB-PROBLEM 10 high: 689709.131427</td>
</tr>
<tr>
<td>GAP = 2361.940101, RELGAP = 0.083787%</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>MASTER PROBLEM 11: 2819338.502316</th>
</tr>
</thead>
<tbody>
<tr>
<td>SUB-PROBLEM 11 high: 692406.351318</td>
</tr>
<tr>
<td>GAP = 2361.940101, RELGAP = 0.083776%</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>MASTER PROBLEM 12: 2819524.204253</th>
</tr>
</thead>
<tbody>
<tr>
<td>SUB-PROBLEM 12 high: 690478.286312</td>
</tr>
<tr>
<td>GAP = 541.528304, RELGAP = 0.019206%</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>MASTER PROBLEM 13: 2819736.994159</th>
</tr>
</thead>
<tbody>
<tr>
<td>GAP = -0.000000, RELGAP = -0.000000%</td>
</tr>
</tbody>
</table>

OPTIMAL SOLUTION FOUND

Expected Cost = 2819736.994159
Decomposition: Observations

Loops can iterate over sets

- Solve a subproblem for each scenario
  - for {s in SCEN} { ... }

One model can represent all subproblems

- Assign loop index s to set S, then solve
  - let S := s;
    solve Sub;

Related solution values can be returned

- Use dual ray to generate infeasibility cuts
  - if Sub.result = "infeasible" then { ... 
    let {i in WHSE}
      supply_price[i,s,nCUT] := Supply[i].dunbdd;
    let {j in STOR}
      demand_price[j,s,nCUT] := Demand[j].dunbdd;
  }
Concluding Observations

Scripts in practice

- Large and complicated
  - Multiple files
  - Hundreds of statements
  - Millions of statements executed
- Run within broader applications

Prospective improvements

- Faster loops
- True script functions
  - Arguments and return values
  - Local sets & parameters
- More database connections
- IDE for debugging
- APIs for popular languages (C++, Java, C#, VB, Python)