



Bounds from Slopes

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WHY BOUND COMPUTATIONS?

- Global optimization
- Assess effects of input uncertainty
- Design – avoid failures

DMG work done at Sandia National Laboratories in Albuquerque, NM, building on work of others cited in the paper:

<http://www.sandia.gov/~dmgay/bounds10.pdf>

Original motivation for this work: assess utility of “uncertainty compiler”.

COMPUTATION STRATEGIES

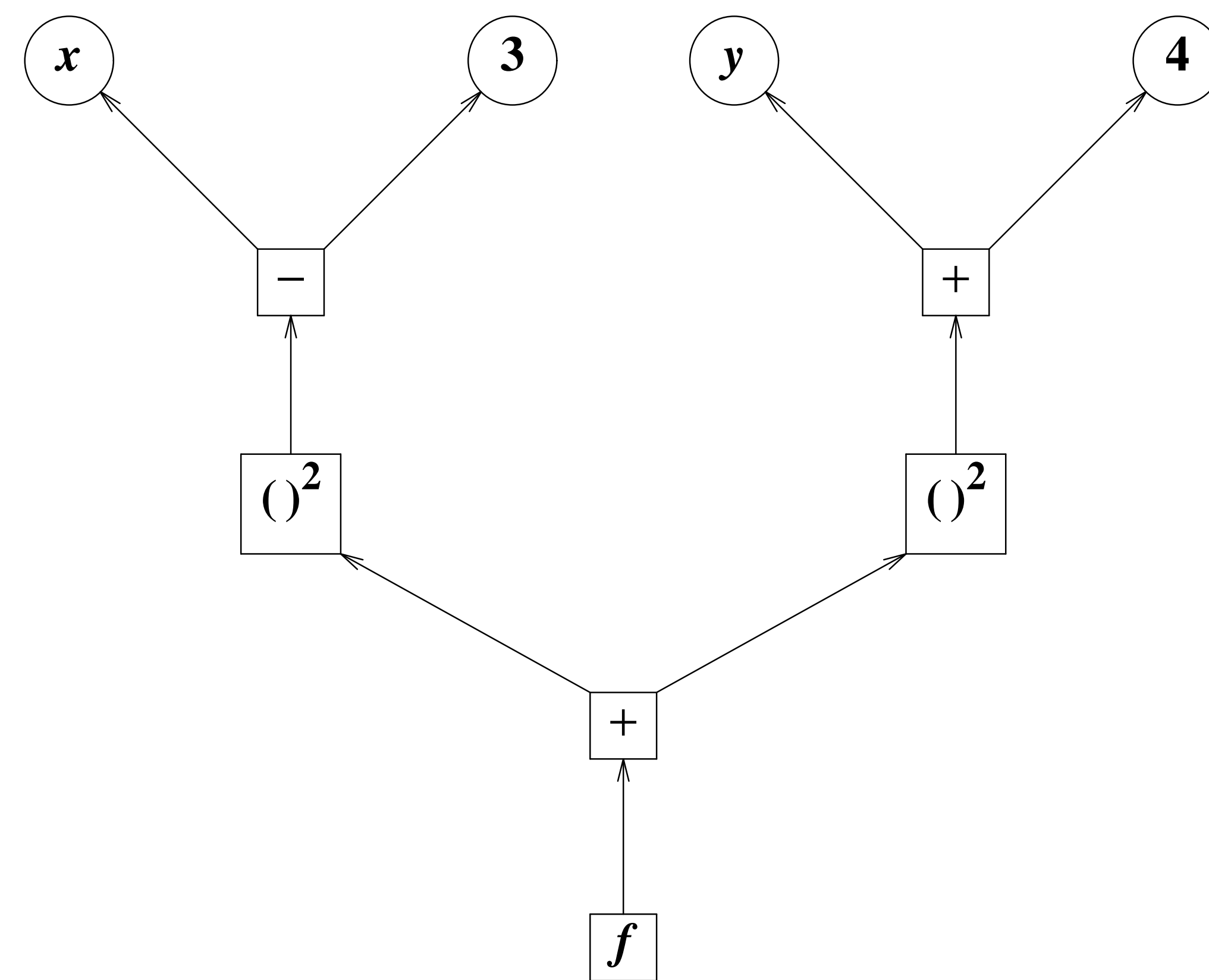
- Compile, link, execute
 - Only choice for large scale?
 - Interpret
 - E.g., *python, awk, perl, ...*
 - Mixture
 - E.g., *Java, AMPL, ...*
- All involve expression graphs.

EXPRESSION GRAPHS

Uses include

- Simplifications before evaluations
- Evaluations
- Derivative computations (AD)
- Bound computations, e.g.,
 - interval
 - Taylor series
 - slope
- Convexity detection.

EXPRESSION GRAPH EXAMPLE



Expression graph for $f = (x-3)^2 + (y+4)^2$

AD-LIKE GRAPH WALKS

Expression-graph walks similar to forward AD can

- compute interval bounds (by interval arithmetic);
- propagate Taylor series;
- compute interval slopes.

SLOPES AND INTERVALS

“Slopes” are divided differences:

$$f[x, z] = \begin{cases} (f(x) - f(z)) / (x - z) & \text{if } x \neq z \\ f'(x) & \text{if } x = z \end{cases}$$

Given interval X , interval evaluation of $f[X, z]$ gives $f[x, z] \in f[X, z] \forall x \in X$.

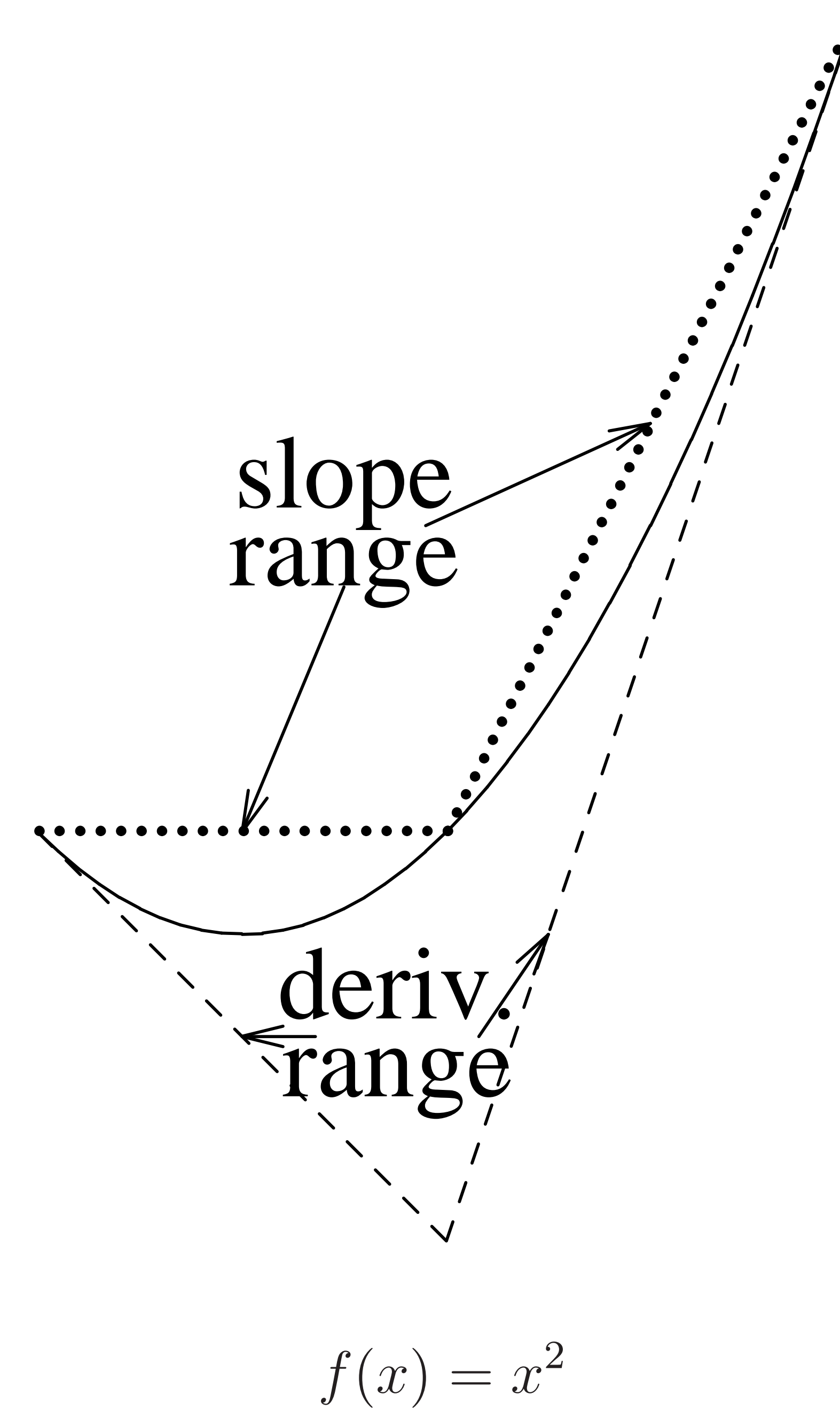
Since $f(x) = f(z) + f[x, z](x - z)$,

$$f(X) \subseteq F_z(X) \doteq f(z) + f[X, z](X - z).$$

Quadratic approximation:

$$\begin{aligned} & \text{width}(F_z(X)) - \text{width}(f(X)) \\ & \leq O(\text{width}(X)^2). \end{aligned}$$

DERIVATIVE VS. SLOPE RANGES



SLOPE ARITHMETIC

Slope arithmetic, analogous to forward AD [Krawczyk & Neumaier, 1985]:

$$\begin{array}{l|l} f = \dots & \Rightarrow f[X, z] = \dots \\ \hline c \in \mathbb{R} & 0 \\ X & 1 \\ g \pm h & g[X, z] \pm h[X, z] \\ g \cdot h & g[X, z] \cdot h(X) + g(z) \cdot h[X, z] \\ g/h & (g[X, z] - h[X, z] \cdot f(z))/h(X) \end{array}$$

Interval slope computations extend readily to n variables and can be done by walk of expression graph.

$$\text{work}(f[X, z]) = O(n \cdot \text{work}(f(x))).$$

Component-wise computation [Hanson, 1968] can be tighter (but slower).

SECOND-ORDER SLOPES

Second-order slopes:

- Scalar: $f[x, z, z] = (f[x, z] - f'(z))/(x - z)$.
- $f(x) = f(z) + f'(z)(x - z) + f[x, z, z](x - z)^2$.
- Often get tighter bounds than with slopes.
- Still just quadratic approximation.
- For $x \in \mathbb{R}^n$, component-wise evaluations have $\text{work}(f[X, z, z]) = O(n \cdot \text{work}(f(x)))$.

SOME BOUNDING TECHNIQUES

interval	$F(X) \supset f(X)$
Taylor 1	$f(z) + F'(X)(X - z)$
slope 1	$f(z) + F[X, z](X - z)$
slope 2	$f(z) + f'(z)(X - z) + F[X, z, z](X - z)^2$
slope 2*	$\{f(z) + f'(z)h + F[X, z, z]h^2 : h \in X - z\}$

Only “slope 2*” and the implementation are new in the present work.

BOUND WIDTHS ON 2 EXAMPLES

Method	Barnes	Sn525
interval	162.417	0.7226
Taylor 1	9.350	0.3609
slope 1	6.453	0.3529
slope 2	3.007	0.1140
slope 2*	2.993	0.1003
true	2.330	0.0903

For problem details, see the paper.

Implementation (not finished) built on facilities in ASL (AMPL/Solver interface library).

Expression graphs from AMPL; use directed roundings (part of IEEE arithmetic standard).

Good for easy development, quick tests.

Contact dmg@acm.org for more details.